# L11 Problem Set 2

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# 1 $LD(S) \wedge LD(T) \wedge LK(T) \implies LD(S \times T)$

Let  $(s_1, t_1), (s_2, t_2), (s_3, t_3) \in S \times T$ . As in the case proved in the lecture slides, the terms to be equated are  $t_{\text{rhs}}$  and  $t_1 \otimes_T t_{\text{lhs}}$ , where

$$(\cdot, t_{\text{lhs}}) = (s_2, t_2) \vec{\oplus} (s_3, t_3), \text{ and}$$
  
 $(\cdot, t_{\text{rhs}}) = (s_1 \otimes_S s_2, t_1 \otimes_T t_2) \vec{\oplus} (s_1 \otimes_S s_3, t_1 \otimes_T t_3)$ 

There are four subcases to consider.

- s<sub>2</sub> = s<sub>2</sub> ⊕<sub>S</sub> s<sub>3</sub> = s<sub>3</sub>
   The proof for the first subcase proved in the slides also works here, since it did not make use of LC(S).
- $s_2 = s_2 \oplus_S s_3 \neq s_3$ Here  $t_{\text{lhs}} = t_2$  and  $s_1 \otimes_S s_2 = s_1 \otimes_S (s_2 \oplus_S s_3)$ . Thus, using LD(S), we have  $s_1 \otimes_S s_2 = (s_1 \otimes_S s_2) \oplus_S (s_1 \otimes_S s_3) =_? s_1 \otimes_S s_3$ . If the  $=_?$  is  $\neq$ , then  $t_{\text{rhs}} = t_1 \otimes_T t_2 = t_1 \otimes_T t_{\text{lhs}}$  as required. If the  $=_?$  is = then  $t_{\text{rhs}} = (t_1 \otimes_T t_2) \oplus_T (t_1 \otimes_T t_3)$ . By LD(T),  $t_{\text{rhs}} = t_1 \otimes_T (t_2 \oplus_T t_3)$ , and by LK(T) this equals  $t_1 \otimes_T t_{\text{lhs}}$  as required.
- $s_2 \neq s_2 \oplus_S s_3 = s_3$ Here  $t_{\text{lhs}} = t_3$  and  $s_1 \otimes_S s_3 = s_1 \otimes_S (s_2 \otimes_S s_3)$ . Thus, using LD(S), we have  $s_1 \otimes_S s_2 = s_1 \otimes_S s_2 \otimes_S s_3 \otimes_S s_3$
- $s_2 \neq s_2 \oplus_S s_3 \neq s_3$ Here  $t_{\text{lhs}} = \bar{0}_T$ . The four possibilities for  $t_{\text{rhs}}$  are  $(t_1 \otimes_T t_2) \oplus_T (t_1 \otimes_T t_3)$ ,  $t_1 \otimes_T t_2$ ,  $t_1 \otimes_T t_3$ , and  $\bar{0}_T$ . The first case, using LD(T), can be written  $t_1 \otimes_T (t_2 \oplus_T t_3)$ . Thus in all cases,  $t_{\text{rhs}}$  is either  $\bar{0}_T$  or  $t_1 \otimes_T t$  for some t. But by LK(T), we know  $t_1 \otimes_T t = t_1 \otimes_T t_{\text{lhs}} = \bar{0}_T$ , thus all cases give the required result.

Note that we assume  $\bar{0}_T$  is an annihilator for  $\otimes_T$ . The theorem should include an extra hypothesis, that if  $\oplus_S$  is not selective (that is, this subcase may occur) then  $\bar{0}_T$  exists and is an annihilator for  $\otimes_T$ .

# 2 $\operatorname{NLD}(S) \vee \operatorname{NLD}(T) \vee (\operatorname{NLC}(S) \wedge \operatorname{NLK}(T)) \implies \operatorname{NLD}(S \times T)$

## **2.1** NLD(S)

From NLD(S) we get witnesses a, b, and c such that  $c \otimes_S (a \oplus_S b) \neq (c \otimes_S a) \oplus_S (c \otimes_S b)$ . We can construct witnesses for NLD( $S \times T$ ) as follows. Let  $t \in T$ . Consider  $(c, t) \otimes ((a, t) \oplus (b, t))$ ,

which equals  $(c,t) \otimes (a \oplus_S b,t') = (c \otimes_S (a \oplus_S b),t'')$  for some t' and t''. Now consider  $((c,t) \otimes (a,t)) \vec{\oplus} ((c,t) \otimes (b,t))$ , which equals  $(c \otimes_S a,t \otimes_T t) \vec{\oplus} (c \otimes_S b,t \otimes_T t) = ((c \otimes_S a) \oplus_S (c \otimes_S b),t''')$  for some t'''. If these were equal, then their first components would be equal. So we would have  $c \otimes_S (a \oplus_S b) = (c \otimes_S a) \oplus_S (c \otimes_S b)$ . But we already know a,b, and c do not satisfy this property. Therefore (a,t),(b,t), and (c,t), for any t, serve as witnesses for  $\text{NLD}(S \times T)$ .

Note we assumed T is not empty, and this condition should be added to the hypotheses of the theorem.

#### $\mathbf{2.2}$ NLD(T)

From NLD(T) we get witnesses a, b, and c such that  $c \otimes_T (a \oplus_T b) \neq (c \otimes_T a) \oplus_T (c \otimes_T b)$ . We can construct witnesses for NLD( $S \times T$ ) as follows. Let  $s \in S$ . Consider  $(s, c) \otimes ((s, a) \overrightarrow{\oplus} (s, b))$ , which equals  $(s, c) \otimes (s \oplus_S s, a \oplus_T b) = (s \otimes_S (s \oplus_S s), c \otimes_T (a \oplus_T b))$ , as long as  $s \oplus_S s = s$ . Now consider  $((s, c) \otimes (s, a)) \overrightarrow{\oplus} ((s, c) \otimes (s, b))$ , which equals  $(s \otimes_S s, c \otimes_T a) \overrightarrow{\oplus} (s \otimes_S s, c \otimes_T b) = ((s \otimes_S s) \oplus_S (s \otimes_S s), (c \otimes_T a) \oplus_T (c \otimes_T b))$ , as long as  $(s \otimes_S s) \oplus_S (s \otimes_S s) = s \otimes_S s$ . If these were equal, then their second components would be equal. So we would have  $c \otimes_T (a \oplus_T b) = (c \otimes_T a) \oplus_T (c \otimes_S b)$ . But we already know a, b, and c do not satisfy this property. Therefore (s, a), (s, b), (s, c), (s, c) for any suitable s, serve as witnesses for NLD( $S \times T$ ).

Similarly to the previous case, we assumed S is not empty. We also need to assume that  $\bigoplus_S$  satisfies  $x \bigoplus_S x = x$  for both x = s and  $x = s \otimes_S s$  for some  $s \in S$ . This would be true if  $\bigoplus_S$  were selective<sup>1</sup> or idempotent, or if  $\bar{0}_S \in S$ .

#### **2.3** $NLC(S) \wedge NLK(T)$

From NLC(S) we get witnesses a, b, and c such that  $c \otimes_S a = c \otimes_S b$  and  $a \neq b$ . From NLK(T) we get witnesses d, e, and f such that  $f \otimes_T d \neq f \otimes_T e$ . We can construct witnesses for NLD( $S \times T$ ) as follows. We proceed by cases on  $a \oplus_S b$  and  $(f \otimes_T d) \oplus_T (f \otimes_T e)$ .

•  $a \oplus_S b = b \wedge (f \otimes_T d) \oplus_T (f \otimes_T e) = f \otimes_T d$  or  $a \oplus_S b = a \wedge (f \otimes_T d) \oplus_T (f \otimes_T e) = f \otimes_T e$ 

<sup>&</sup>lt;sup>1</sup>selectivity implies idempotence